

Exam Logical Verification

January 14, 2008

(6) exercises.

be given in Dutch or English. Good luck!

This exercise is concerned with first-order propositional logic
 simply typed λ -calculus ($\lambda \rightarrow$).

Proof in prop1 showing that the following formula is a tautology:

$$((A \rightarrow B \rightarrow A) \rightarrow A) \rightarrow A$$

type-derivation in $\lambda \rightarrow$ corresponding to the proof in 1a.

the following simply typed λ -terms:

$\lambda x : ?. \lambda y : ?. \lambda z : ?. x z y$
 $\lambda x : ?. \lambda y : ?. \lambda z : ?. x (z y)$
 $\lambda x : ?. \lambda y : ?. y x x$

This exercise is concerned with first-order predicate logic (pred1)
 with dependent types (λP).

Proof in pred1 showing that the following formula is a tautology:

$$(\forall x. P(x) \rightarrow Q(x)) \rightarrow (\forall x. P(x)) \rightarrow (\forall x. Q(x))$$

λP -term corresponding to the formula in 2a.

Find an inhabitant in λP of the answer to 2b.

representing the propositions is declared
`prop : Set.`
 The operation `on prop` is a binary operator *)
`imp : prop -> prop -> prop.`
 Use infix notation `=>` for `imp` *)
`" := imp (right associativity, at level 100)`
 It expresses if a proposition in `prop` is valid
`p`) is inhabited then `p` is valid
`p`) is not inhabited then `p` is not valid
`r T : prop -> Prop.`

types of the parameters `imp_introduction` and
`imp_elimination` rule and elimination rule of the `imp`

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This exercise is concerned with Coq.

definition of an inductive datatype `two` with ex
 s)

the induction principle for the datatype `two`.

s)

for the following inductive definition of the pre

`inductive even : nat -> Prop :=`
 `| 0 : even 0`
 `| nSS : forall n:nat , even n -> even (S n)`

if possible, an inhabitant of the following:

`even 0`
`even 1`
`even 2`

oints)

a fixed point definition of a function that takes
 gives as output the sum of its elements (an
 ty).

oints)

Final note is (the total amount of points plus 1

