

Exam Logical Verification

May 31, 2013

There are six (6) exercises.

Answers may be given in Dutch or English. Good luck!

Exercise 1. (*5+5+6 points*)

This exercise is concerned with first-order propositional logic (**prop1**) and simply typed λ -calculus ($\lambda \rightarrow$).

- a) Show that the following formula is a tautology of minimal **prop1**:
 $(A \rightarrow A \rightarrow B) \rightarrow (C \rightarrow A) \rightarrow (C \rightarrow B)$.
- b) Give the type derivation in $\lambda \rightarrow$ corresponding to the proof of 1a.
- c) Give, if possible, closed inhabitants in $\lambda \rightarrow$ of the following types:
 $((B \rightarrow A \rightarrow B) \rightarrow A) \rightarrow A$
 $A \rightarrow A \rightarrow B$
 $A \rightarrow B \rightarrow A$

Exercise 2. (*5+3+5+3 points*)

This exercise is concerned with first-order predicate logic (**pred1**) and λ -calculus with dependent types (λP).

- a) Show that the following formula is a tautology of minimal **pred1**:
 $(\forall x. (P(x) \rightarrow Q(x))) \rightarrow \forall x. ((Q(x) \rightarrow R(x)) \rightarrow P(x) \rightarrow R(x))$.
- b) Give a λP -term corresponding to the formula in 2a.
(Use **Terms** for the domain that is quantified over.)
- c) Give a closed inhabitant in λP of the answer to 2b.
- d) Consider the following question Q :
is the formula $\forall x. P(x) \rightarrow (\forall y. P(y) \rightarrow A) \rightarrow A$ a tautology?
What is the counterpart of the question Q in λ -calculus/type theory?

Exercise 3. (5+3+5+3 points)

This exercise is concerned with second-order propositional logic (**prop2**) and polymorphic λ -calculus (**$\lambda 2$**).

- a) Show that the following formula is a tautology of minimal **prop2**:
 $a \rightarrow \forall b. ((\forall c. a \rightarrow c) \rightarrow b)$.
- b) Give the $\lambda 2$ -term corresponding to the formula in 3a.
- c) Give a closed inhabitant in $\lambda 2$ of the answer to 3b.
- d) What is the proof checking problem? Is it decidable for $\lambda 2$?

Exercise 4. (2+4+4+6 points)

This question is concerned with various typing issues.

- a) The typing rules we considered use $*$ and \square .
 What is/are the counterpart(s) in Coq of $*$ and of \square ?
- b) We define **and** $C D$ with $C : *$ and $D : *$ in $\lambda 2$ as follows:

$$\text{and } C D := \Pi a : *. (C \rightarrow D \rightarrow a) \rightarrow a$$

Assume an inhabitant $P : \text{and } C D$. Give an inhabitant of C (provide the informal typing derivation).

- c) Give the polymorphic identity in $\lambda 2$.
 Next, assume $\text{nat} : *$ and show how the polymorphic identity is instantiated to the identity on **nat** using application and β -reduction.
- d) Consider the typing rules for a product $\Pi x:A. B$:

$$\text{product } (\lambda \rightarrow \text{and } \lambda P \text{ and } \lambda 2) \quad \frac{\Gamma \vdash A : * \quad \Gamma, x : A \vdash B : *}{\Gamma \vdash \Pi x:A. B : *}$$

$$\text{product } (\lambda P) \quad \frac{\Gamma \vdash A : * \quad \Gamma, x : A \vdash B : \square}{\Gamma \vdash \Pi x:A. B : \square}$$

$$\text{product } (\lambda 2) \quad \frac{\Gamma \vdash A : \square \quad \Gamma, x : A \vdash B : *}{\Gamma \vdash \Pi x:A. B : *}$$

We assume $\text{nat} : *$ and $\text{vec} : \text{nat} \rightarrow *$. Explain informally how the following products can be typed using the appropriate product rule.

- (i) $\text{nat} \rightarrow \text{nat}$
- (ii) $\Pi a : *. a \rightarrow a$
- (iii) $\Pi n : \text{nat}. \text{vec } n$

Exercise 5. (*5+4+4 points*)

This exercise is concerned with inductive datatypes in Coq.

- a) Give the definition of an inductive datatype **three** with exactly three elements. Also, give the type of **three_ind** for the induction principle on **three_ind**.
- b) Give the inductive definition of the datatype **natsnoclist** of lists of natural numbers, but where the constructor for adding an element to a list adds this element at the end.
- c) Give the type of **natsnoclist_ind** for the induction principle on **natsnoclist**.

Exercise 6. (*4+5+4 points*)

This exercise is concerned with inductive predicates in Coq.

- a) Consider the inductive predicate for less-than-equal in Coq:

```
Inductive le (n:nat) : nat -> Prop :=  
| le_n : le n n  
| le_S : forall m:nat , le n m -> le n (S m) .
```

Prove that $1 \leq 2$, that is, give an inhabitant of `le (S 0) (S (S 0))`.

- b) Give the definition of an inductive predicate **evenlist** on the usual datatype **natlist** (of lists of natural numbers) such that **evenlist** `l` holds exactly if the list `l` has an even number of elements.
- c) Complete the following definition of conjunction in Coq:

```
Inductive and (A : Prop) (B : Prop) : Prop :=
```

The note for the exam is (the total amount of points plus 10) divided by 10.