

This exam consists of two pages

- 1a MINIX3 has adopted the **client-server model** to structure itself. Explain what this model entails, and notably what the role of the kernel is. 5pt
- 1b Drivers in MINIX3 run as ordinary user-space processes. In what sense does this impose restrictions for communicating with hardware controllers, and how is that solved? 5pt
- 1c Sketch the flow of control when a user-space process calls the library routine `read(fd,buffer,bytes)` in a monolithic operating system. Hint: use a diagram. 5pt
- 1d In many cases, the hardware offers support for multiple rings of protection for programs. How can we make use of this support when organizing an operating system? 5pt
- 2a The semantics of the *atomic swap* machine instruction is defined as follows. Show how this instruction can be used to protect a critical section. 5pt
- `swap(inout boolean a, inout boolean b){ temp = a; a = b; b = temp;}`
- 2b Consider the program on page 2, which is to be executed as a separate MINIX3 user-space **lock manager** process. The core of the program is formed by the functions `do_down()` and `do_up()` which are standard operations on counting semaphores. Given `lock_manager()`, give a pseudo-code implementation of the function `do_down(sema)`. 5pt
- 2c Also give a sketch of the implementation of `do_up()`. 5pt
- 2d Returning `SUSPEND` by `do_down()` has the result of suspending a process. Explain which process that is, and how this blocking is actually effectuated. 5pt
- 3a Explain how MINIX3 (and many other operating systems) simulate multiple timers using a single clock. Draw a figure to explain your answer. 5pt
- 3b Explain the difference between **character devices** and **block devices**, and why making this distinction can be helpful for improving I/O. Hint: think of writing a stream of bytes to disk. 5pt
- 4a Explain the principle working of the `fork()` system call. 5pt
- 4b **Copy-on-write** is a technique by which a block of memory is filled with data from a specific source only when first written to. How can this technique help in optimizing the implementation of `fork()`? Be precise! 5pt
- 5a Explain what the `mount()` system call does by means of an example. Explain your example! 5pt
- 5b `Mount()` changes fields in inodes and in-memory copies of superblocks. Explain these changes. 5pt
- 5c Consider the following operations that are carried out on a formatted, but otherwise empty USB stick. Explain what the result will be when listing the directory contents (by means of the last operation `ls`). 5pt
- | | |
|--|--------------------------------------|
| <code>mount /dev/sdb1 /usbstick</code> | <i>Mount the USB stick</i> |
| <code>cd /usbstick</code> | <i>Enter the directory</i> |
| <code>mkdir test</code> | <i>Create a subdirectory 'test.'</i> |
| <code>touch test/x</code> | <i>Create a file 'x' in 'test.'</i> |
| <code>mount /dev/sdb1 test</code> | <i>Mount the USB stick again</i> |
| <code>ls test</code> | <i>List the directory contents</i> |
- 5d Explain precisely what happened with the superblock table and inode table after the two mount operations from the previous example have been carried out. 5pt
- 6a What is a protection domain? 5pt

6b Give a practical example of how to switch from one protection domain to another, and explain how such a switch could be implemented by an operating system.

5pt

```
01 PUBLIC int lock_manager(){
02   int result, s, proc_nr;
03   struct mproc *rmp;
04   while (TRUE) {
05     receive(ANY, &msg_in);
06     who = msg_in.m_source;          /* who sent the message      */
07     sema = msg_in.m5_l1;           /* which semaphore is this? */
08     call_request = msg_in.m5_i1;    /* which operation is requested? */
09     mp = &mproc[who];
10     switch(call_request){
11       DOWN: result = do_down(sema); break;
12       UP:   result = do_up(sema); break;
13     }
14
15     /* Send the results back to the user to indicate completion. */
16     if (result != SUSPEND) setreply(who, result); /* Prepare reply message */
17     /* Send out all pending reply messages, including the answer to
18      * the call just made above.
19      */
20
21     for (proc_nr = 0, rmp = mproc; proc_nr < NR_PROCS; proc_nr++, rmp++) {
22       if ((rmp->mp_flags & REPLY) == REPLY ){
23         send(proc_nr, &rmp->mp_reply);
24         rmp->mp_flags &= ~REPLY;
25       }
26     }
27   }
28   return(OK);
29 }
```

<p>Grading: The final grade is calculated by adding the scores per question (maximum: 90 points), and adding 10 bonus points. The maximum total is therefore 100 points.</p>
