This exam has 6 pages and 8 exercises.

The result will be computed as (total number of points plus 10) divided by 10.

Answers may be given in either English or Dutch.

Please motivate all answers!

## 1. Equivalence relations (7 + 8 points)

On the set  $A := \{1, 2, 3, 4\} \times \{1, 2, 3, 4\}$ , consider the relation R that is defined by the description

$$\langle n_1, k_1 \rangle R \langle n_2, k_2 \rangle \iff |n_1 - k_1| = |n_2 - k_2|$$

- (a) Show that R is an equivalence relation.
- (b) Explicitly write down all the different equivalence classes of the equivalence relation R on A, and give a complete system of representatives.

Solution:

- (a) We must show that R is reflexive, symmetric, and transitive.
  - (i) Since  $|n-k|=|n-k|, \langle n,k\rangle R\langle n,k\rangle$  for all  $\langle n,k\rangle\in A$ . Hence, R is reflexive.
  - (ii) Suppose  $\langle n_1, k_1 \rangle R \langle n_2, k_2 \rangle$ . Then  $|n_1 k_1| = |n_2 k_2|$ , hence  $|n_2 k_2| = |n_1 k_1|$ , and therefore  $\langle n_2, k_2 \rangle R \langle n_1, k_1 \rangle$ . So R is symmetric.
  - (iii) Suppose  $\langle n_1, k_1 \rangle R \langle n_2, k_2 \rangle$  and  $\langle n_2, k_2 \rangle R \langle n_3, k_3 \rangle$ . Then  $|n_1 k_1| = |n_2 k_2|$  and  $|n_2 k_2| = |n_3 k_3|$ , which implies that  $|n_1 k_1| = |n_3 k_3|$ , and therefore,  $\langle n_1, k_1 \rangle R \langle n_3, k_3 \rangle$ . So R is transitive.
- (b) The equivalence classes are

$$\{\langle 1, 1 \rangle, \langle 2, 2 \rangle, \langle 3, 3 \rangle, \langle 4, 4 \rangle \}$$

$$\{\langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 2, 3 \rangle, \langle 3, 2 \rangle, \langle 3, 4 \rangle, \langle 4, 3 \rangle \}$$

$$\{\langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 2, 4 \rangle, \langle 4, 2 \rangle \}$$

$$\{\langle 1, 4 \rangle, \langle 4, 1 \rangle \}$$

Therefore, a complete system of representatives is

$$\{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle\}$$

### 2. Injections and Surjections (4 + 4 points)

For each of the two functions defined below, determine whether or not the function is injective, and whether or not the function is surjective, and explain your answers.

- (a) The function  $f: \mathbb{Z} \to \mathbb{N}$  defined by f(n) = |n-2|.
- (b) The function  $g: CSstudents \to N$  that assigns to every Computer Science student at VU Amsterdam his or her student number.

Solution:

- (a) f is not injective because, for example, f(1) = |-1| = f(3) = |1| = 1. f is surjective because for every  $k \in \mathbb{N}$ , f(k+2) = |k| = k.
- (b) Since different Computer Science students have different student numbers, g is injective. g is not surjective, because the number of Computer Science students is finite, while the codomain of g is infinite.

### 3. Function composition (4 + 4 points)

We are given the functions  $exp_2 \colon R \to R$ ,  $inv \colon R \to R$  and  $inc \colon R \to R$  defined by

$$exp_2(x) = 2^x$$
  $inv(x) = \frac{1}{x}$   $inc(x) = x + 1$ 

(a) Give an explicit description (i.e. a formula for f(x)) of the function f defined by

$$f:=\exp_2\circ inv\circ inc.$$

(b) Express the function  $g: \mathbb{R} \to \mathbb{R}$  defined by

$$g(x) := \frac{1}{2^{x-2}}$$

as a composition of the functions  $exp_2$ , inv and inc and their inverses.

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Solution:

(a) 
$$f(x) = 2^{\frac{1}{x+1}}$$

(b)  $g = inv \circ exp_2 \circ inc^{-1} \circ inc^{-1}$ 

### 4. Induction (4 + 10 points)

Consider the sequence  $(t_n)_{n=1}^{\infty}$  of numbers defined recursively by

$$t_1 := 0,$$
  $t_{n+1} := t_n + 2(n+1).$ 

- (a) Calculate the terms  $t_2$ ,  $t_3$ ,  $t_4$  and  $t_5$  of this sequence.
- (b) Prove by mathematical induction that for all  $n \geq 1$ ,

$$t_n = (n+2)(n-1).$$

Solution:

- (a)  $t_2 = t_1 + 4 = 4$ ,  $t_3 = t_2 + 6 = 10$ ,  $t_4 = t_3 + 8 = 18$ ,  $t_5 = t_4 + 10 = 28$ .
- (b) First,  $t_1 = 0$  by definition, but also  $(1+2)(1-1) = 3 \cdot 0 = 0$ . Therefore,  $t_n = (n+2)(n-1)$  holds for n = 1. This proves the base case of the induction.

Now for the inductive step. Assume that  $t_m = (m+2)(m-1)$ , where  $m \ge 1$  is arbitrary (this is our inductive hypothesis). By the recursive definition and the inductive hypothesis,

$$t_{m+1} = t_m + 2(m+1)$$
 (by definition)  
=  $(m+2)(m-1) + 2(m+1)$  (by the inductive hypothesis)  
=  $m^2 + 2m - m - 2 + 2m + 2$   
=  $m^2 + 3m$ 

On the other hand,

$$(m+1+2)(m+1-1) = (m+3)m = m^2 + 3m$$

This shows that if the inductive hypothesis is true, then also

$$t_{m+1} = (m+1+2)(m+1-1)$$

is true. This completes the inductive proof.

#### 5. Axioms for semantic equivalence (12 points)

Show using the axioms for semantic equivalence that

$$\phi \wedge (\psi \vee (\neg \phi \wedge \chi)) \equiv \phi \wedge \psi$$

Solution:

$$\phi \wedge (\psi \vee (\neg \phi \wedge \chi)) \equiv (\phi \wedge \psi) \vee (\phi \wedge (\neg \phi \wedge \chi)) \quad \text{(distributivity)}$$

$$\equiv (\phi \wedge \psi) \vee ((\phi \wedge \neg \phi) \wedge \chi) \quad \text{(associativity)}$$

$$\equiv (\phi \wedge \psi) \vee (\bot \wedge \chi) \quad \text{(complement)}$$

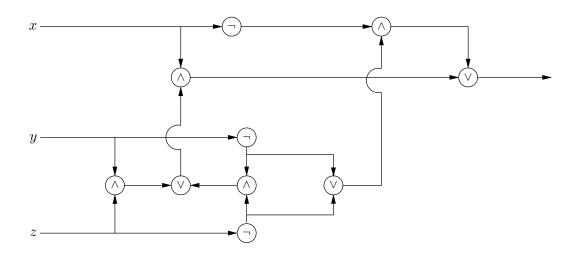
$$\equiv (\phi \wedge \psi) \vee (\chi \wedge \bot) \quad \text{(commutativity)}$$

$$\equiv (\phi \wedge \psi) \vee \bot \quad \text{(domination)}$$

$$\equiv \phi \wedge \psi \quad \text{(identity)}$$

# **6. Logic circuit** (8 points)

Give the propositional formula that corresponds to the following logic circuit:



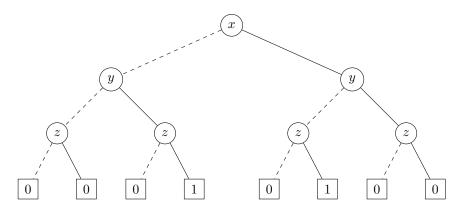
Solution:  $(x \land ((y \land z) \lor (\neg y \land \neg z))) \lor (\neg x \land (\neg y \lor \neg z))$ 

# **7. OBDD** (4 + 8 points)

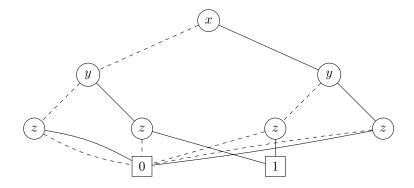
- (a) Represent  $(x \oplus y) \land z$  by means of a binary decision tree, with respect to the variable ordering x, y, z.
- (b) Reduce this binary decision tree to an ordered binary decision diagram.

Solution:

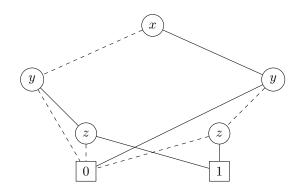
(a)



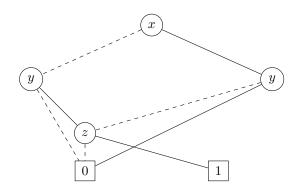
(b) First apply C1 to collapse leaves:



Next apply C2 twice to remove the most left-hand and the most right-hand z nodes:



Finally apply C3 to collapse the two remaining z nodes:



# 8. Predicate logic (3 + 3 + 7 points)

Suppose there is a collection of dolls and there are two boxes. Consider the following three unary predicates:

- D(x): x is a doll

-  $B_1(x)$ : x is in the first box

-  $B_2(x)$ : x is in the second box

- (a) Express the phrase "Every doll is in one of the two boxes" in predicate logic.
- (b) Express the phrase "All dolls are in the first box or all dolls are in the second box" in predicate logic.
- (c) Are the formulas in (a) and (b) semantically equivalent? If so, argue they coincide on all models. If not, give a model on which they produce different truth values.

#### Solution:

(a) This sentence can either be interpreted as an inclusive or

$$\forall x (D(x) \rightarrow B_1(x) \lor B_2(x))$$

or as an exclusive or

$$\forall x (D(x) \to B_1(x) \oplus B_2(x)).$$

(b) Again this sentence can either be interpreted as an inclusive or

$$\forall x (D(x) \to B_1(x)) \lor \forall x (D(x) \to B_2(x))$$

or as an exclusive or

$$\forall x (D(x) \to B_1(x)) \oplus \forall x (D(x) \to B_2(x)).$$

(c) The two formulas are not semantically equivalent.

Consider a model with a set of two elements a and b, where D(a), D(b),  $B_1(a)$ , and  $B_2(b)$ . (But  $B_1(b)$  and  $B_2(a)$  don't hold.)

The formula of part (a) holds in this model, while the formula of part (b) does not.