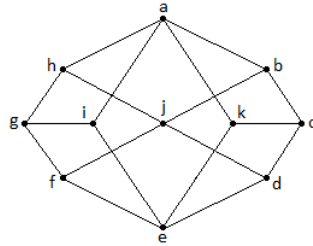
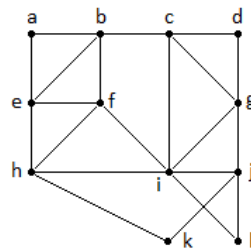


- Correct answers without full explanation score badly, so justify your steps.
- The use of books, lecture notes, calculators, etc. is *not* allowed.

(1) Consider the following graph  $G$ :



- Is  $G$  bipartite? Either give the bipartite sets, or prove that such a partition does not exist.
  - Is  $G$  Hamiltonian? Either give a Hamiltonian cycle, or prove that it does not exist.
  - Is  $G$  planar? Either give a planar representation, or prove that it does not exist.
  - Determine the maximum number of edges in a matching in  $G$ . (As always, explain/verify your assertion.)
- (2) Let  $G$  be a graph of order  $n$  with less than  $n(n-1)/2$  edges. Show that  $G$  is  $(n-1)$ -colorable.
- (3) Let  $G$  be a graph with adjacency matrix  $A$ . show that:  
 $G$  is bipartite  $\Leftrightarrow$  for all positive odd integers  $k$  we have  $\text{Tr}(A^k) = 0$ .
- (4) Let  $G$  be a  $k$ -regular Eulerian graph of even order. Prove that  $G$  has an even number of edges.
- (5) Draw a labeled tree whose Prüfer sequence is 1, 3, 1, 5, 1, 7. Describe your steps.
- (6) Use Hierholzer's algorithm to find an Eulerian circuit in the following graph. Use  $R_1 : a, b, c, g, i, h, e, a$  as your initial circuit. Describe your steps.



- (7) Let  $G$  be a connected planar graph that is not a tree, with  $v$  vertices,  $e$  edges, and  $f$  regions/faces (in a planar representation). Suppose that  $G$  has at least  $m$  vertices of degree 1.
- Show that  $2e \geq 3f + 2m$ .  
 (Hint: consider  $\sum_R (b(R) + 2c(R))$ , where  $c(R)$  denotes the number of edges that come into contact with region  $R$  but not with another region.)
  - Show that  $3(v-2) \geq e + 2m$ .  
 (Note: you are allowed to use the inequality given at (a) above for answering (b).)