Databases

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Databases

A database (DB) is a collection of data with

- a certain logical structure
- a specific semantics
- a specific group of users

A database management system (DBMS) allows to

- create, modify and manipulate a database
- query (retrieve) the data using a query language
- support persistent storage of large amounts of data
- enable durability and recovery from failure
- control access to the data by many users in parallel
 - without unexpected interactions among users (isolation)
 - actions on the data should never be partial (atomicity)

Why not just store data in files?

Why not just store data in files?

- no query language
- weak logical structure (limited to directories)
- no efficient access
 - searching through a large file can take hours
- no or limited protection from data loss
- no access control for parallel manipulation of data

So we need database management systems...

Motivation for database management systems

data independence

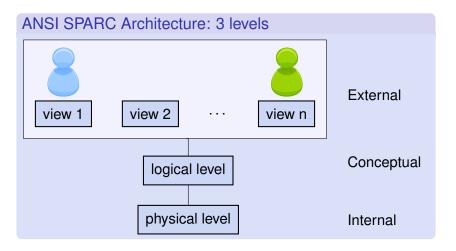
- logical view on the data independent of physical storage
- user interacts with a simple view on the data
- behind the scenes (invisible for the user) are complex storage structures that allow rapid access and manipulation

avoidance of duplication

- different views on the same database
 - for different users or different applications
 - hiding parts of the data for privacy or security

This is achieved by the ANSI SPARC Architecture . . .

View of Data



- Different applications might use different views
- Data is stored only once at the physical level
 - good for consistency

View of Data

ANSI SPARC Architecture: 3 levels

- View level:
 - application programs hide details of data types
 - hide information (e.g. exam grade) for privacy or security
- Logical level: also called 'conceptual schema'
 - describes data stored in the database, and
 - relations among the data
- Physical level:
 - how the data is stored
 - disk pages, index structures, byte layout, record order

This ensures logical and physical data independence...

Data Independence

Logical data independence

Logical data independence is the ability to modify the logical schema without breaking existing applications

applications access the views, not the logical database

Physical data independence

Physical data independence is the ability to modify the physical schema without changing the logical schema

- e.g. a change in workload might cause the need for
 - different indexing structures
 - different database engine
 - distributing the database on multiple machines
 - **...**

Relational Model

In this course, we work with **relational databases**. Their view and logical level represented data as **relations/tables**.

Example relational database instance customer account id city depositor name street accountnr 1928374 Johnson 12 Alma Palo Alto 1928374 101343 3211231 Jones 34 Main Harisson 3211231 217343 0192837 Smith 7 South Rye 0192837 201762 row = tuple record: (3211231, Jones, 34 Main, Harisson)

In the **pure relational model**, a table is a **set** of tuples:

- has no duplicate tuples (rows)
- no order on the tuples

Motivation for database management systems

- high-level declarative query languages
 - query tells what you want, independent of storage structure
 - efficient data access (automatic query optimisation)

Declarative Query Languages

Queries should:

- describe what information is sought
- not prescribe how to retrieve the desired information

Relational databases usually use SQL as query language . . .

Imperative vs. Declarative Languages

Kowalski

Imperative/procedural languages:

- explicit control
- implicit logic

Declarative/non-procedural languages:

- implicit control
- explicit logic

Examples of declarative languages

- logic programming (e.g. Prolog),
- functional programming (e.g. Haskell),
- markup languages (e.g. HTML), . . .

SQL = Structure Query Language

SQL is a declarative data manipulation language. The user describes conditions the requested data is required to fulfil.

SQL Query

```
SELECT ID
FROM CUSTOMER
WHERE NAME = 'Jones' AND CITY = 'Harisson'
```

More concise than imperative languages:

- less expensive program development
- easier maintenance

Database system will optimise the query:

decides how to execute the query as fast as possible Users (usually) do not need to think about efficiency.

Motivation for database management systems

- well-defined data models & data integrity constraints
 - relational model
 - meta language for describing
 - data
 - data relationships
 - data constraints

SQL can be used for table and constraint definitions ...

Relational Model: Schema

Database schema

= structure of the database = relations + constraints

Example schema

- customer(<u>id</u>, name, street, city)
- account(depositor → customer, accountnr)

Database instance

= actual content ('state') of the database at some moment

Example instance

customer			
<u>id</u>	name	street	city
1928374	Johnson	12 Alma	Palo Alto
0192837	Smith	4 North	Rye

account			
depositor	accountnr		
1928374	101343		
0192837	215569		

Integrity Constraints

Example schema with key constraints

- customer(<u>id</u>, name, street, city)
 Primary key constraint on <u>id</u>
- account(depositor → customer, accountnr) Foreign key constraint on depositor

Various types of constraints:

- data types, constrained data types (domains)
- columns constraints (e.g. unique, nullability, counter, ...)
- check constraints (logical expression for domain integrity)
 (e.g. age >= 18 AND age <= 150)

SQL DDL (Data Definition Language)

Creating a table with constraints

```
CREATE TABLE solved (
  id INT AUTO_INCREMENT,
  name VARCHAR(40) NOT NULL,
  homework NUMERIC(2) NOT NULL,
  points NUMERIC(2) NOT NULL CHECK (points <= 10),
  PRIMARY KEY (id)
);</pre>
```

Note the data types and constraints!

```
solved

id
name
homework
points
```

Creating a view

```
CREATE VIEW solved_homework AS

SELECT id, name, homework FROM solved;
```

How to Design Database Schemes?

How to design database schemes?

- Entity-relationship models (ER models)
- UML class diagrams

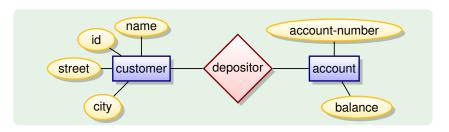
They are

- widely used for database design
- usually converted to the relational model

Entity Relationship Model

Entity relationship model

- entities = objects
 - e.g. customers, accounts, bank branches
- relationship between entities
 - e.g. account 101343 is held by customer Johnson
 - relationship set descriptor links customers with accounts

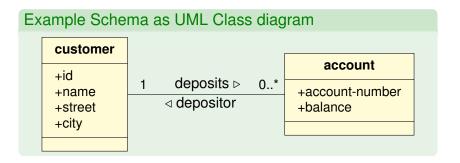


UML Class Diagram

UML class diagrams

- frequently used in database design
- similar to E/R diagrams:

 $Entities/Relationships \implies Classes/Associations$



Motivation for database management systems

- multiple users, concurrent access
 - transactions with ACID properties

A **transaction** is a collection of operations that performs a single logical function in a database application.

Database management system ensures ACID properties

- Atomicity: transaction executes fully (commit) or not at all (abort)
- Consistency: database remains in a consistent state where all integrity constraints hold
- Isolation: multiple users can modify the database at the same time but will not see each others partial actions
- Durability: once a transaction is committed successfully, the modified data is persistent, regardless of disk crashes

Summary

Why Database Management Systems?

- data independence
 - logical view on the data independent of physical storage
- avoidance of duplication
 - different views on the same database
- high-level declarative query languages (what, not how)
 - efficient data access, automatic query optimisation
- data models & data integrity (consistency)
- multiple users, concurrent access
 - transactions with ACID properties
- persistent storage, safety and high availability
 - safety against failure (backup/restore)
- scalability (data could by much larger than main memory)
 - indexing, scalable algorithms
- security