Resit Data Structures and Algorithms 2018-2019

Tuesday January 8, 2019, 18.30-21.15

6 exercises

#### Shortly explain your answers unless otherwise specified



## Exercise 1. (5+5+5 points)

This exercise is concerned with sorting.

- (a) Give the worst-case time complexity in terms of  $\mathcal{O}$  for sorting an array of n numbers using (i) selection sort, (ii) quicksort, (iii) merge sort, (iv) counting sort, (v) heapsort. No motivation needed.
- (b) Give the recursion tree for the application of merge sort to the input array [1, 5, 6, 4, 8, 7, 2, 3].
- (c) Give a recurrence equation describing the worst-case time complexity of merge sort.

Solve your recurrence equation to give the worst-case time complexity of merge sort in terms of  $\Theta$ .

### Exercise 2. (6+6+3 points)

This exercise is concerned with heaps.

- (a) Consider the max-heap H = [9, 8, 6, 4, 7, 5, 1, 2, 3]. Apply 'on the fly' the algorithm HeapExtractMax (that removes and returns the maximum element of a max-heap) to H; give your answer in pictures.
- (b) Apply 'on the fly' the algorithm for bottom-up max-heap construction to the array A = [1, 2, 3, 4, 5, 6, 7]; give your answer in pictures.
- (c) What is the worst-case time complexity in terms of  $\mathcal{O}$  for the algorithm for bottum-up max-heap construction on an input-array of length n?

#### Exercise 3. (6+4 points)

(a) Give an implementation (use pseudo-code) of a queue with operations enqueue and dequeue, using a singly linked list, where we have the following: for an element x in the list, we have operations x.next and x.key with the suggested meaning. For a list L we have operation L.head.

Give (no motivation needed) the worst-case time complexity of your operations.

(b) Does an algorithm A with worst-case time complexity in  $\mathcal{O}(n)$  always perform better than an algorithm B with worst-case time complexity in  $\mathcal{O}(n^2)$ ? (Explain your answer.)

## Exercise 4. (5+5+5+5 points)

This exercise is concerned with binary search trees and AVL-trees.

- (a) Give in pictures all binary search trees with labels 1, 2, 3.
- (b) What is the worst-case time complexity, in terms of  $\mathcal{O}$ , of adding an element to (i) a Binary Search Tree (BST) consisting of n elements, (ii) a min-heap consisting of n elements, (iii) a AVL-tree consisting of n elements? No motivation needed.
- (c) Consider binary search trees implemented as a linked structure (for example, v.key is the key at node v).

Give pseudo-code for a non-recursive procedure for searching a key in a a sub-tree of a binary search tree. The inputs for the procedure are a pointer x to a node in the binary search tree, and a key k. The output is a pointer to a node with key k if such node exists in the sub-tree rooted at node x, and nil otherwise.

(d) Construct an AVL-tree by inserting one by one the keys

### $3 \quad 5 \quad 4 \quad 2 \quad 1 \quad 6 \quad 7$

starting from the empty tree. After each insertion, rebalance the tree if needed. Give your answer in pictures.

# Exercise 5. (5+5+5 points)

This exercise is concerned with greedy choice.

- (a) Consider the knapsack01 optimization problem: given a set S of items all with benefit and weight, and given a maximum weight W, select an optimal choice  $C \subseteq S$  in the sense that the total benefit of C is maximal under the constraint that the total weight of C does not exceed W.
  - Give a small example showing that the greedy choice for an item with maximal benefit does not necessarily yield an optimal solution.
- (b) Consider the activity selection problem: given a finite set of activities all with start time and finish time, select a maximum-size subset of mutually compatible activities.

Argue that the greedy choice for an activity with minimal end time yields an optimal solution.

(c) We consider the following problem: we are given a finite set of activities, each with start time and finish time. We have available infinitely many lecture rooms.

Describe (informally but precisely) a greedy algorithm for scheduling all activities using a minimum number of lecture rooms.

# Exercise 6. (5+4+6 points)

(a) Consider the algorithm for a longest common subsequence (LCS) of input sequences  $X = \langle x_1, \dots, x_m \rangle$  and  $Y = \langle y_1, \dots, y_n \rangle$ :

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Algorithm LCS(X,Y):

new array C[0 \dots m, 0 \dots n]

for i := 0 to m do

C[i,0] := 0

for j := 0 to n do

C[0,j] := 0

for i := 1 to m do

for j := 1 to n do

if x_i = y_j then

C[i,j] := C[i-1,j-1] + 1

else

C[i,j] := max(C[i,j-1], C[i-1,j])

return C
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Apply the algorithm to the following input:  $X = \langle A, B, C, D, E, F \rangle$  and  $Y = \langle A, B, D, F, C, D \rangle$ . Give your answer in the form of a table and give explicitly the longest common subsequence(s) that is (are) found.

- (b) What is the worst-case time complexity of the LCS algorithm of 6(a) in terms of  $\mathcal{O}$ ? (Explain your answer.)
- (c) Describe an algorithm that takes as input an array of integers and that finds a longest increasing subsequence of the input-array. What is the worst-case time complexity of your algorithm?