

Midterm Exam Data Structures and Algorithms 2015-2016

Friday September 25, 2015, 15.30-17.00

7 exercises

Answers may be given in English or in Dutch.



**Exercise 1.** (*5+5+5 points*)

This exercise is concerned with  $\mathcal{O}$ .

- (a) Consider the pseudo-code for the procedure **partition** for quick-sort. It works on a sub-array  $A[p \dots r]$  of an array  $A$  of length  $n$ . Explain why its worst-case running time is in  $\mathcal{O}(n)$ .

```
Algorithm partition( $A, p, r$ ):  
   $x := A[r]$   
   $i := p - 1$   
  for  $j = p$  to  $r - 1$  do  
    if  $A[j] \leq x$  then  
       $i := i + 1$   
      exchange  $A[i]$  with  $A[j]$   
  exchange  $A[i + 1]$  with  $A[r]$   
  return  $i + 1$ 
```

- (b) What is the worst-case time complexity in terms of  $\mathcal{O}$  of the following loop? Motivate your answer.

```
Algorithm LoopA( $n$ ):  
   $a := 0$   
  for  $i := 1$  to  $2n$  do  
    for  $j := 1$  to  $i$  do  
       $a := a + i$ 
```

- (c) Give an example of an algorithm that has linear best-case time complexity, and quadratic worst-case time complexity. Motivate your answer.

**Exercise 2.** (5+5 points)

This exercise is concerned with sorting.

- (a) What is the worst-case time complexity in terms of  $\mathcal{O}$  of selection sort, bubble sort, merge sort, quick sort, counting sort. You do not need to motivate your answer.
- (b) What is the lower bound for the worst-case time complexity of comparison-based sorting? You do not need to motivate your answer.

**Exercise 3.** (5+6 points)

This exercise is concerned with insertion sort; its pseudo-code is given:

**Algorithm** insertionSort( $A, n$ ):

```
for  $j := 2$  to  $n$  do
     $key := A[j]$ 
     $i := j - 1$ 
    while  $i \geq 1$  and  $A[i] > key$  do
         $A[i + 1] := A[i]$ 
         $i := i - 1$ 
     $A[i + 1] := key$ 
```

- (a) Apply insertion sort to the array  $[3, 4, 1, 5, 2]$ .  
Give (at least) the intermediate result after every iteration of the for-loop.
- (b) Give pseudo-code for a recursive version of insertion sort, that sorts the first  $n - 1$  elements recursively, and inserts the last element at the right position.

**Exercise 4.** (5+5+5 points)

- (a) Give the merge-sort tree for sorting the input array  $[3, 6, 1, 7, 5, 8, 2, 4]$ .
- (b) Solve the following recurrence equation for merge sort  
(you may assume  $n$  to be a power of 2):

$$T(n) = \begin{cases} 1 & \text{als } n = 1 \\ 2T(\frac{n}{2}) + n & \text{als } n > 1 \end{cases}$$

- (c) Explain the divide-and-conquer programming paradigm.

**Exercise 5.** (6+7 points)

This exercise is concerned with singly linked lists. In a node  $v$  we have operations  $v.next$  and  $v.element$  with the suggested meaning. For a list  $L$  we have operations  $L.first$  and  $L.last$  with the suggested meaning. Do not assume predefined operations on lists.

- (a) Implement a queue with the operations `enqueue` and `dequeue` as a singly linked list.
- (b) Give a  $\Theta(n)$ -time non-recursive procedure that reverses a singly-linked list of  $n$  elements, only using a constant amount of additional storage.

**Exercise 6.** (5+5 points)

This exercise is concerned with hashing.

- (a) What is the worst-case time complexity of searching for hashing where collisions are solved using chaining? Briefly explain your answer.
- (b) We consider a hash table of length 11 (an array with indices  $0 \dots 10$ ), and the hash function  $h(k) = k \bmod 11$ . Add the following numbers in this order to the initially empty hash table:

1   13   2   24   10   12

solving collisions by open addressing with linear probing.

**Exercise 7.** (6+5+5 points)

This exercise is concerned with binary search trees and AVL-trees.

- (a) Give pseudo-code for a recursive procedure for searching a key in a binary search tree. The inputs for the procedure are a pointer  $x$  to a node, and a key  $k$ . The output is a pointer to a node with key  $k$  if such node exists, and `nil` otherwise.
- (b) What is the worst-case time complexity of search in a binary search tree? And in an AVL tree? Briefly motivate your answers.
- (c) Construct an AVL-tree by inserting the numbers

3   2   1   5   4   6

one by one, starting from the empty tree. After each insertion, rebalance the tree if needed. Give your answer in pictures (with comments if needed).

*The mark for the midterm is (the total number of points plus 10) divided by 10.*